

# Indexed hierarchies and internal structure in the ordinal class in intuitionistic set theories

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# Ordinals and ordinal-indexed hierarchies

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Fix a set-like partial order  $\mathcal{R}$  on the class  $\text{Tr}$  of all transitive sets, which is no coarser than  $\subseteq$  yet no finer than  $\subseteq$ , i.e.

$$x = y \vee x \in y \implies x \mathcal{R} y; \quad x \mathcal{R} y \implies x \subseteq y,$$

and we consider the  $\mathcal{R}$ -predecessor set

$$P_{\mathcal{R}}(x) = \{y : y \mathcal{R} x\}.$$

E.g. when  $\mathcal{R} = \subseteq$ ,  $P_{\mathcal{R}}(x) = \mathcal{P}(x)$ ; when  $\mathcal{R} = \subseteq$ ,  $P_{\mathcal{R}}(x) = x^+ = x \cup \{x\}$ .  
We can also get the first-order definable subsets  $\text{Def}(x)$ , etc.

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$$\begin{aligned}H_{\mathcal{R}}(\emptyset) &= \emptyset, \\H_{\mathcal{R}}(\alpha^+) &= P_{\mathcal{R}}(H_{\mathcal{R}}(\alpha)), \\H_{\mathcal{R}}(\lambda) &= \bigcup_{\alpha \in \lambda} H_{\mathcal{R}}(\alpha) \quad \text{for limit } \lambda.\end{aligned}$$

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E.g.  $V_{\alpha}$ ,  $L_{\alpha}$ , etc.

Alternatively, for any set  $x$ , we can define

$$H_{\mathcal{R}}(x) = \bigcup_{y \in x} P_{\mathcal{R}}(H_{\mathcal{R}}(y)).$$

One easily verifies that this gives the same results on ordinals.

# Ordinals and ordinal-indexed hierarchies

Fix a set-like partial order  $\mathcal{R}$  on the class  $\text{Tr}$  of all transitive sets, which is no coarser than  $\subseteq$  yet no finer than  $\in$ .

## Question

Is there a maximal class  $O_{\mathcal{R}} \subseteq V$  on which  $H_{\mathcal{R}}$  is an injection?

'Maximal' means for any  $x \in V$ , there exists  $y \in O_{\mathcal{R}}$  such that  $H_{\mathcal{R}}(x) = H_{\mathcal{R}}(y)$ .

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**Answer (ZF).** Always take  $O_{\mathcal{R}} = \text{Ord}$ .

Injectivity follows from linearity of ordinals (note  $x \in y$  implies  $H_{\mathcal{R}}(x) \in H_{\mathcal{R}}(x^+) \subseteq H_{\mathcal{R}}(y)$ ).

For any  $x \in V$ , simply observe  $H_{\mathcal{R}}(x) = H_{\mathcal{R}}(\text{rank}(x))$ .

# Intuitionistic set theory

Intuitionistic Zermelo–Fraenke set theory IZF is:

ZF, minus Excluded Middle, with Axiom of Foundation/Regularity replaced by Set Induction:

$$\forall x (\forall y \in x \varphi(y) \rightarrow \varphi(x)) \rightarrow \forall x \varphi(x).$$

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## Wellknown folklore

IZF is consistent with Ord being non-linear, or formally

$$\exists \alpha, \beta \in \text{Ord} (\beta \subseteq \alpha \wedge \beta \notin \alpha \wedge \beta \neq \alpha).$$

This is easily implementable in multiple types of ‘intuitionistic models’, e.g. models on Kripke frames, realisability models or Heyting-valued models.

# Intuitionistic set theory

## Corollary

IZF is consistent with the class function  $\alpha \mapsto V_\alpha$  being non-injective.

**Proof.** Take  $\alpha, \beta \in \text{Ord}$  such that  $\beta \subseteq \alpha$ ,  $\beta \notin \alpha$  and  $\beta \neq \alpha$ . Consider

$$V_{\alpha^+} = \mathcal{P}(V_\alpha).$$

But  $V_\beta \subseteq V_\alpha$ , i.e.  $\mathcal{P}(V_\beta) \subseteq \mathcal{P}(V_\alpha)$ , thus

$$V_{\alpha^+ \cup \beta^+} = \mathcal{P}(V_\alpha) \cup \mathcal{P}(V_\beta) = V_{\alpha^+}.$$

Clearly  $\beta \notin \alpha^+$ , so  $\alpha^+ \neq \alpha^+ \cup \beta^+$ . □

# The intuition

## Question

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It is still reasonable to expect  $O_{\mathcal{R}} \subseteq \text{Ord}$ . Ordinals in  $O_{\mathcal{R}}$  should satisfy some additional 'closure properties'.

# The intuition

## Question

Is there a maximal class  $O_{\mathcal{R}} \subseteq V$  on which  $H_{\mathcal{R}}$  is an injection?

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**Example (Paul Taylor, 1996):** We say an ordinal  $\alpha \in \text{Ord}$  is *plump* if

- every  $\beta \in \alpha$  is plump;
- for every  $\gamma \in \text{Ord}$ , if  $\gamma$  is plump and  $\gamma \subseteq \beta$  for some  $\beta \in \alpha$ , then  $\gamma \in \alpha$ .

As it turns out, the class of plump ordinals is precisely a maximal class on which  $\alpha \mapsto V_{\alpha}$  is injective.

## The intuition

A similar idea comes from Robert Lubarsky, *Intuitionistic L* (1993):

### Theorem

Although it is open whether  $\text{IZF} \vdash \text{Ord} \subseteq L$ , we can avoid it and still prove  $\text{IZF} \vdash (V = L)^L$ .

**Proof.** For any  $\alpha \in \text{Ord}$ , we construct  $\alpha^* \in L$  as

$$\alpha^* = \left\{ \beta \in \text{Ord} : L_{(\alpha_{\text{aug}} + \text{H}n)^-} \models \text{def}(L_\beta) \subseteq L_\alpha \right\},$$

such that  $L_\alpha = L_{\alpha^*} \in L$ . □

Note that  $\text{def}(L_\beta) \subseteq L_\alpha$  can be equivalently written as  $L_\beta \in L_\alpha$ , since if  $L_\beta \in \text{def}(L_\gamma)$  for some  $\gamma \in \alpha$ , then also

$$\text{def}(L_\beta) \subseteq \text{def}(L_\gamma) \subseteq L_\alpha.$$

# The formal construction

We now directly define

$$O_{\mathcal{R}} = \{\alpha \in \text{Ord} : \forall \beta \in \alpha^+ \forall \gamma \in \text{Ord} \\ (H_{\mathcal{R}}(\gamma) \in H_{\mathcal{R}}(\beta) \rightarrow \exists \gamma' \in \beta H_{\mathcal{R}}(\gamma') = H_{\mathcal{R}}(\gamma))\}.$$

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## Proposition

$H_{\mathcal{R}}$  is injective on  $O_{\mathcal{R}}$ .

**Proof.** Use induction. Suppose  $\beta, \beta' \in O_{\mathcal{R}}$  and  $H_{\mathcal{R}}(\beta) = H_{\mathcal{R}}(\beta')$ , we want to show  $\beta = \beta'$ . For  $\gamma \in \beta$ ,

$$H_{\mathcal{R}}(\gamma) \in H_{\mathcal{R}}(\beta) = H_{\mathcal{R}}(\beta'),$$

i.e.  $H_{\mathcal{R}}(\gamma) = H_{\mathcal{R}}(\gamma')$  for  $\gamma' \in \beta'$ . By inductive hypothesis,  $\gamma = \gamma' \in \beta'$ .

The converse is symmetrical. □

## $\mathcal{R}$ -successors

Are there many elements in  $O_{\mathcal{R}}$ ?

We want to first find the  $\mathcal{R}$ -successor of  $\alpha \in O_{\mathcal{R}}$ , i.e. some  $\alpha^{\mathcal{R}+} \in O_{\mathcal{R}}$  such that

$$H_{\mathcal{R}}(\alpha^{\mathcal{R}+}) = H_{\mathcal{R}}(\alpha^+) = P_{\mathcal{R}}(H_{\mathcal{R}}(\alpha)).$$

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We consider the set

$$s = \{\{\gamma \in \alpha : H_{\mathcal{R}}(\gamma) \in x\} : x \in P_{\mathcal{R}}(H_{\mathcal{R}}(\alpha))\}.$$

For any  $\beta \in \text{Ord}$ , if  $H_{\mathcal{R}}(\beta) \in P_{\mathcal{R}}(H_{\mathcal{R}}(\alpha))$ , then for each  $\gamma \in \beta$ , there is  $\gamma' \in \alpha$  such that  $H_{\mathcal{R}}(\gamma') = H_{\mathcal{R}}(\gamma)$ . Thus

$$H_{\mathcal{R}}(\beta) = \bigcup_{\gamma \in \beta} P_{\mathcal{R}}(H_{\mathcal{R}}(\gamma)) = \bigcup_{\gamma' \in t} P_{\mathcal{R}}(H_{\mathcal{R}}(\gamma')) = H_{\mathcal{R}}(t)$$

for  $t = \{\gamma \in \alpha : H_{\mathcal{R}}(\gamma) \in H_{\mathcal{R}}(\beta)\} \in s$ . Easy to verify  $t \in O_{\mathcal{R}}$ .

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it suffices to define  $\alpha^{\mathcal{R}+} = \{\beta \in s : \beta \in O_{\mathcal{R}}, H_{\mathcal{R}}(\beta) \mathcal{R} H_{\mathcal{R}}(\alpha)\}$ .

### Lemma

*For any  $\alpha \in O_{\mathcal{R}}$ ,  $\alpha^{\mathcal{R}+} \in O_{\mathcal{R}}$  is the least element in  $O_{\mathcal{R}}$  that contains  $\alpha$ .*

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### Lemma

For  $s \subseteq O_{\mathcal{R}}$ ,  $\sup s = \bigcup s \in O_{\mathcal{R}}$ .

So we can define  $\text{rk}_{\mathcal{R}} : V \rightarrow O_{\mathcal{R}}$  as

$$\text{rk}_{\mathcal{R}}(x) = \bigcup_{y \in x} \text{rk}_{\mathcal{R}}(y)^{\mathcal{R}^+}.$$

# Maximality of $O_{\mathcal{R}}$

$$\text{rk}_{\mathcal{R}}(x) = \bigcup_{y \in x} \text{rk}_{\mathcal{R}}(y)^{\mathcal{R}^+}.$$

Using  $H_{\mathcal{R}}(\alpha^{\mathcal{R}^+}) = P_{\mathcal{R}}(H_{\mathcal{R}}(\alpha))$ , we prove by induction that:

## Proposition

For any set  $x$ ,

$$H_{\mathcal{R}}(x) = H_{\mathcal{R}}(\text{rk}_{\mathcal{R}}(x)).$$

## Corollary

For any set  $x$ , there exists  $\alpha \in O_{\mathcal{R}}$  such that  $H_{\mathcal{R}}(x) = H_{\mathcal{R}}(\alpha)$ , i.e.  $O_{\mathcal{R}}$  is a maximal class on which  $H_{\mathcal{R}}$  is injective.

# Consistency of class separations

- ZF proves that  $O_{\mathcal{R}} = \text{Ord}$  for all  $\mathcal{R}$ .
- When  $\mathcal{R} = \subseteq$ ,  $H_{\mathcal{R}}(x) = \text{rk}_{\text{Ord}}(x)$ , so we always have  $O_{\subseteq} = \text{Ord}$ .

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- Let  $\mathcal{R}_{\text{def}}$  be (first-order) definable subset relation, so that  $H_{\mathcal{R}_{\text{def}}}(\alpha) = L_{\alpha}$ . It is consistent that  $2 \notin O_{\subseteq}$  (when 1 has undecidable subsets, i.e.  $\alpha \subseteq 1$  with  $\emptyset \in \alpha \vee \emptyset \notin \alpha$  failing to hold), but  $2 = L_2 \in O_{\mathcal{R}_{\text{def}}}$  always, so consistently  $O_{\subseteq} \neq O_{\mathcal{R}_{\text{def}}}$ .

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- Using realisability arguments, we can verify

Theorem (W., upcoming PhD thesis)

IZF is consistent with  $\alpha \neq \beta \wedge L_{\alpha} = L_{\beta}$ .

Thus also consistently  $O_{\mathcal{R}_{\text{def}}} \neq \text{Ord} = O_{\subseteq}$ .

## Plump ordinals are useful!

It remains open from Lubarsky (1993) whether  $\text{IZF} \vdash \text{Ord} \subseteq L$ . However,

### Theorem (W., 2026)

$O_{\subseteq} \subseteq L$ . More specifically, for any  $\alpha \in O_{\subseteq}$ ,

$$\alpha = \{\beta \in L_{\alpha} : L_{\alpha} \models \beta \in O_{\subseteq}\}.$$

(We need to phrase  $O_{\subseteq}$  as Taylor's plump ordinals, to avoid mentioning the objects  $H_{\subseteq}(\beta) = V_{\beta}$  in  $L_{\alpha}$ .)

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Combined with the use of incomparable ordinals, it follows that

### Theorem (W., 2026)

Fix any set  $x$  in any ZFC universe, then there is a Heyting algebra  $\mathbb{H}$  such that the Heyting extension satisfies

$$V^{\mathbb{H}} \models \mathcal{P}(\check{x}) \in L.$$

# Ordinal arithmetic

Fix any  $\mathcal{R}$ , we can let

$$\alpha +_{\mathcal{R}} \beta = \alpha \cup \bigcup_{\gamma \in \beta} (\alpha +_{\mathcal{R}} \gamma)^{\mathcal{R}+}, \quad \alpha \cdot_{\mathcal{R}} \beta = \bigcup_{\gamma \in \beta} (\alpha \cdot_{\mathcal{R}} \gamma +_{\mathcal{R}} \alpha).$$

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## Lemma

*For any  $\alpha, \beta, \gamma \in \text{Ord} = \mathcal{O}_{\subseteq}$ , if  $\alpha + \beta = \alpha + \gamma$ , then  $\beta = \gamma$ ; if  $\alpha + \beta \in \alpha + \gamma$ , then  $\beta \in \gamma$ .*

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These proofs are different (and indeed, the  $\mathcal{O}_{\subseteq}$  one is even simpler), and I do not know any uniform way to do this for arbitrary  $\mathcal{R}$ .

# Thank you!

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- [1] Robert Seth Lubarsky, *Intuitionistic L*, Logical methods: In honor of Anil Nerode's sixtieth birthday (John Newsome Crossley, Jeffrey Brian Remmel, Richard Arnold Shore, and Moss Eisenberg Sweedler, eds.), Birkhäuser Boston, 1993, pp. 555–571.
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